Multigrid Methods for the Biharmonic Problem with Cahn-Hilliard Boundary Conditions

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1 Introduction

Let $\Omega \subset \mathbb{R}^2$ be a bounded polygonal domain, $V = \{v \in H^2(\Omega) : \frac{\partial v}{\partial n} = 0 \text{ on } 12 \partial \Omega\}$ and $f \in L_2(\Omega)$. In this paper we consider multigrid methods for the following 13 biharmonic problem: Find $u \in V$ such that 14

$$\int_{\Omega} \nabla^2 u : \nabla^2 v \, dx = \int_{\Omega} f v \, dx \qquad \forall v \in V, \tag{1}$$

where $\nabla^2 w: \nabla^2 v = \sum_{i,j=1}^2 w_{x_i x_j} v_{x_i x_j}$ is the inner product of the Hessian matrices of w_{15} and v. Under the (assumed) compatibility condition, 16

$$\int_{\Omega} f \, dx = 0,\tag{2}$$

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the biharmonic problem (1) is solvable and the solution is unique up to an additive 17 constant. Furthermore we have an elliptic regularity estimate 18

 $\|\hat{u}\|_{H^{2+\alpha}(\Omega)} \le C \|f\|_{L_2(\Omega)}$ (3)

for the solution \hat{u} of (1) that satisfies $\int_{\Omega} \hat{u} dx = 0$. Note that, unlike the biharmonic ¹⁹ problem with the boundary conditions of clamped plates, the index of elliptic regu-²⁰ larity α in (3), which is determined by the angles of Ω , can be close to 0 even if Ω ²¹ is convex (cf. [2]).²²

The essential boundary condition $\partial u/\partial n = 0$ and the natural boundary condition ²³ $\partial (\Delta u)/\partial n = 0$ satisfied by the solution *u* of (1) appear in the Cahn-Hilliard model ²⁴ for phase separation phenomena (cf. [8]). In particular, the boundary value problem ²⁵ (1) appears when the Cahn-Hilliard equation is discretized in time by an implicit ²⁶ method and the resulting nonlinear fourth order elliptic boundary value problem is ²⁷ solved by an Newton iteration. ²⁸

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We will describe a C^0 interior penalty method for (1) in Sect. 2 and introduce in ²⁹ Sect. 3 multigrid methods that are based on a new smoother. The convergence prop-³⁰ erties of the multigrid methods are briefly discussed in Sect. 4, followed by numerical ³¹ results in Sect. 5. ³²

2 A Quadratic C⁰ Interior Penalty Method

 C^0 interior penalty methods (cf. [6, 9]) are discontinuous Galerkin methods for fourth ³⁴ order problems. Let \mathscr{T}_h be a simplicial triangulation of Ω , $V_h \subset H^1(\Omega)$ be the associated P_2 Lagrange finite element space (cf. [5]), and \hat{V}_h be the subspace of V_h consisting of functions with zero mean, i.e., $v \in V_h$ belongs to \hat{V}_h if and only if $\int_{\Omega} v \, dx = 0$. The quadratic C^0 interior penalty method for (1) is to find $\hat{u}_h \in \hat{V}_h$ such that

$$a_h(\hat{u}_h, v) = \int_{\Omega} f v \, dx \qquad \forall v \in \hat{V}_h, \tag{4}$$

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where

$$a_{h}(w,v) = \sum_{T \in \mathscr{T}_{h}} \int_{T} \nabla^{2} w : \nabla^{2} v \, dx + \sum_{e \in \mathscr{E}_{h}} \int_{e} \left\{ \left\{ \frac{\partial^{2} w}{\partial n^{2}} \right\} \right\} \left[\left[\frac{\partial v}{\partial n} \right] \right] ds + \sum_{e \in \mathscr{E}_{h}} \int_{e} \left\{ \left\{ \frac{\partial^{2} v}{\partial n^{2}} \right\} \right\} \left[\left[\frac{\partial w}{\partial n} \right] \right] ds + \sum_{e \in \mathscr{E}_{h}} \frac{\sigma}{|e|} \int_{e} \left[\left[\frac{\partial w}{\partial n} \right] \right] \left[\left[\frac{\partial v}{\partial n} \right] \right] ds.$$
(5)

Here \mathscr{E}_h is the set of the edges in \mathscr{T}_h , $\{\{\partial^2 v/\partial n^2\}\}$ (resp. $[\![\partial v/\partial n]\!]$) is the average of 40 the second normal derivative of *v* (resp. the jump of the first normal derivative of *v*) 41 across an edge, |e| is the length of the edge *e*, and $\sigma > 0$ is a penalty parameter. 42

The quadratic C^0 interior penalty method is consistent. It is also stable if σ is 43 sufficiently large, which is assumed to be the case. (The magnitude of σ is related to 44 certain inverse estimates. It can be taken to be 5 in practice.) It can be shown (cf. [3]) 45 that the solution \hat{u}_h of (4) satisfies the following error estimate: 46

$$\|\hat{u} - \hat{u}_h\|_h \le Ch^{\alpha} \|f\|_{L_2(\Omega)},$$
 (6)

where \hat{u} is the zero mean solution of (1), α is the index of elliptic regularity in (3), 47 and the norm $\|\cdot\|_h$ is given by 48

$$\|v\|_{h}^{2} = \sum_{T \in \mathscr{T}_{h}} |v|_{H^{2}(T)}^{2} + \sum_{e \in \mathscr{E}_{h}} |e|^{-1} \| [\![\partial v / \partial n]\!]\|_{L_{2}(e)}^{2}.$$

 C^0 interior penalty methods have certain advantages over other finite element ⁴⁹ methods for fourth order problems. They are simpler than conforming methods ⁵⁰ which require C^1 elements. They come in a natural hierarchy that can capture smooth ⁵¹ solutions efficiently, which is not the case for classical nonconforming methods. Unlike mixed methods they preserve the positive definiteness of the continuous problem ⁵³ and are easier to develop for more complicated problems (cf. [9]). ⁵⁴

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Another significant advantage of C^0 interior penalty methods comes from the 55 fact that the underlying finite element spaces are standard spaces for second order 56 problems. (Note that the essential boundary condition for (1) is only enforced weakly 57 in (4) and the finite element space V_h does not involve any boundary condition.) 58 Therefore multigrid solves for second order problems can be easily implemented as 59 a preconditioner. By using such a preconditioner in the smoothing steps of multigrid 60 algorithms for fourth order problems, the performance of the smoother and hence 61 the overall performance of the multigrid algorithms can be significantly improved. 62 This approach was carried out in [7] for the biharmonic problem with the boundary 63 conditions of clamped plates. Below we will use this approach to develop multigrid 64 methods for (4).

3 Multigrid Methods

Let \mathscr{T}_k ($k = 0, 1, \cdots$) be a sequence of simplicial triangulations obtained from the 67 initial triangulation \mathscr{T}_0 by uniform refinement. We will use V_k (resp. $a_k(\cdot, \cdot)$) to denote 68 the finite element space (resp. the bilinear form for the C^0 interior penalty method) 69 associated with \mathscr{T}_k .

Let V'_k be the dual space of V_k and $\hat{V}_k = \{v \in V_k : \int_{\Omega} v \, dx = 0\}$ be the zero-mean 71 subspace of V_k . We can identify \hat{V}'_k with the subspace of V'_k whose members annihilate 72 the constant functions, i.e., $\hat{V}'_k = \{\gamma \in V'_k : \langle \gamma, 1 \rangle = 0\}$, where $\langle \cdot, \cdot \rangle$ is the canonical 73 bilinear form between a vector space and its dual. 74

Let the operator $A_k : V_k \longrightarrow \hat{V}'_k$ be defined by $\langle A_k v, w \rangle = a_k(v, w)$ for all $v, w \in V_k$. 75 We can then rewrite the discrete problem (4) as $A_k \hat{u}_k = \phi_k$, where $\hat{u}_k \in \hat{V}_k$ and $\phi_k \in \hat{V}'_k$ 76 satisfies $\langle \phi_k, v \rangle = \int_{\Omega} f v \, dx$ for all $v \in V_k$. Below we will develop multigrid algorithms 77 for equations of the form 78

$$A_k z = \psi \tag{7}$$

where $z \in \hat{V}_k$ and $\psi \in \hat{V}'_k$.

There are two ingredients in the design of multigrid algorithms. First of all, we so need intergrid transfer operators to move data between consecutive levels. Since since intergrid transfer operators to move data between consecutive levels. Since I_{k-1}^{k} : so $V_{k-1} \longrightarrow V_k$ to be the natural injection and the fine-to-coarse operator $I_k^{k-1} : V_k' \longrightarrow I_{k-1}^{k}$ to be the transpose of I_{k-1}^{k} with respect to the canonical bilinear forms, i.e., so V_{k-1}^{k-1} to be the transpose of I_{k-1}^{k} with respect to that I_{k-1}^{k} maps \hat{V}_{k-1} into \hat{V}_k and so consequently I_k^{k-1} maps \hat{V}'_k into \hat{V}'_{k-1} .

The second ingredient is a good smoother that can damp out the highly oscillatory 87 part of the error of an approximate solution so that the remaining part of the error 88 can be captured accurately on a coarser grid. Here we take advantage of the fact that 89 the P_2 Lagrange finite element space is a standard space for second order problems 90 to incorporate a multigrid Poisson solve in the smoother. Let $L_k : \hat{V}_k \longrightarrow \hat{V}'_k$ be the 91 discrete Laplace operator defined by 92

$$\langle L_k v, w \rangle = \int_{\Omega} \nabla v \cdot \nabla w \, dx \qquad \forall v, w \in \hat{V}_k.$$
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We take $S_k^{-1}: \hat{V}'_k \longrightarrow \hat{V}_k$ to be an approximate inverse of L_k obtained from a multigrid 94 Poisson solve such that 95

$$\langle S_k v, v \rangle \approx |v|_{H^1(\Omega)}^2 \qquad \forall v \in \hat{V}_k.$$
 (8)

The smoothing step in our multigrid algorithms for (7) is then given by

$$z_{\text{new}} = z_{\text{old}} + \lambda_k S_k^{-1} (\psi - A_k z_{\text{old}}), \tag{9}$$

where λ_k is a damping factor chosen so that the spectral radius $\rho(\lambda_k S_k^{-1} A_k)$ is <2. It 97 follows from (8) and standard inverse estimates (cf. [5]) that we can take $\lambda_k = Ch_k^2$. 98 Note that the computational cost of (9) is proportional to the dimension of \hat{V}_k , which 99 implies that the overall computational costs of the multigrid algorithms in Sects. 3.1 100 and 3.2 are also proportional to the dimension of \hat{V}_k . 101

We can now describe the *V*-cycle and *W*-cycle algorithms (cf. [10]) in terms of $_{102}$ the integrid transfer operators and the smoothing scheme. $_{103}$

3.1 V-Cycle Algorithm

The *V*-cycle algorithm computes an approximate solution $MG_V(k, \psi, z_0, m)$ of (7) 105 with initial guess $z_0 \in \hat{V}_k$ and *m* pre-smoothing and *m* post-smoothing steps. For 106 k = 0, we take $MG_V(0, \psi, z_0, m)$ to be the output of a direct solve. For $k \ge 1$, we 107 compute $MG_V(k, \psi, z_0, m)$ recursively in three steps. 108

Pre-smoothing For $1 \le \ell \le m$, compute z_ℓ recursively by

$$z_{\ell} = z_{\ell-1} + \lambda_k S_k^{-1} (\psi - A_k z_{\ell-1})$$

Coarse Grid Correction Compute

$$z_{m+1} = z_m + I_{k-1}^k MG_V(k-1, \rho_{k-1}, 0, m),$$

where $\rho_{k-1} = I_k^{k-1}(\psi - A_k z_m) \in \hat{V}'_{k-1}$ is the transferred residual of z_m . *Post-smoothing* For $m + 2 \le \ell \le 2m + 1$, compute z_ℓ recursively by 112

$$z_{\ell} = z_{\ell-1} + \lambda_k S_k^{-1} (\psi - A_k z_{\ell-1}).$$

The final output is $MG_V(k, \psi, z_0, m) = z_{2m+1}$.

3.2 W-Cycle Algorithm

The *W*-cycle algorithm computes an approximate solution $MG_W(k, \psi, z_0, m)$ of (7) 115 with initial guess $z_0 \in \hat{V}_k$ and *m* pre-smoothing and *m* post-smoothing steps. The only 116 difference between the *V*-cycle algorithm and the *W*-cycle algorithm is in the coarse 117 grid correction step, where the coarse grid algorithm is applied twice to the coarse 118 grid residual equation. More precisely, we have 119

$$z_{m+\frac{1}{2}} = MG_W(k-1,\rho_{k-1},0,m),$$

$$z_{m+1} = z_m + MG_W(k-1,\rho_{k-1},z_{m+\frac{1}{2}},m).$$

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Remark 1. For simplicity we have described the multigrid algorithms in terms of the 121 space \hat{V}_k where the bilinear form $a_k(\cdot, \cdot)$ is nonsingular. But the multigrid Poisson 122 solve S_k^{-1} (and hence the *V*-cycle and *W*-cycle algorithms) can be implemented on 123 V_k for $k \ge 1$. The implementation of multigrid algorithms for the singular Neumann 124 problem is discussed for example in [1]. 125

4 Convergence Properties

Let $z_0 \in \hat{V}_k$ be the initial guess and $z_{\dagger} \in \hat{V}_k$ be the output of the *V*-cycle or *W*-cycle 127 algorithm for (7). Numerical results indicate that 128

$$||z - z_{\dagger}||_{a_h} \le Cm^{-\alpha} ||z - z_0||_{a_h}, \tag{10}$$

where α is the index of elliptic regularity in (3) and $\|\cdot\|_{a_h} = \sqrt{a_h(\cdot, \cdot)}$ is the energy 129 norm, provided that the number of smoothing steps $m \ge m_*$. Here m_* is a sufficiently 130 large positive integer independent of k. In particular the multigrid algorithms are 131 contractions for sufficiently large m and the contraction numbers are bounded away 132 from 1 uniformly. A similar estimate was obtained in [7] for the boundary conditions 134 will be carried out in [4] where general fourth order problems are considered. 135

A significant benefit of including a multigrid Poisson solve in the smoothing step 136 (9) is that the resulting smoothing property is similar to that for second order prob-137 lems (cf. [7]) so that the contraction number estimate (10) is also similar to that for 138 second order problems. Indeed, because of the estimate (8), we can derive a smooth-139 ing property for (9) with respect to a family of mesh dependent norms $\||\cdot|\|_{s,k}$ such 140 that $\||\cdot|\|_{0,k} \approx |\cdot|_{H^1(\Omega)}$ and $\||\cdot|\|_{1,k} \approx |\cdot|_{H^2(\Omega)}$ on the space \hat{V}_k . Note that the smoothing 141 properties of standard smoothers for second order problems are described in terms 142 of mesh dependent norms $\||\cdot|\|_{s,k}$ such that $\||\cdot|\|_{0,k} \approx \|\cdot\|_{L_2(\Omega)}$ and $\||\cdot|\|_{1,k} \approx |\cdot|_{H^1(\Omega)}$ 143 on the finite element spaces. The good performance of the smoothing step (9) is due 144 to the similarity between the Hilbert scales $[H^1(\Omega), H^2(\Omega)]$ and $[L_2(\Omega), H^1(\Omega)]$. 145

If we use a standard smoother such as the Richardson relaxation in a multigrid 146 algorithm for (7), then the smoothing property will be determined by the Hilbert 147 scale $[L_2(\Omega), H^2(\Omega)]$. In this case the estimate (10) will be replaced by the estimate 148

$$\|z - z_{\dagger}\|_{a_h} \le Cm^{-\alpha/2} \|z - z_0\|_{a_h},\tag{11}$$

which means that the effect of 100 smoothing steps without the preconditioner is 149 roughly equivalent to the effect of 10 smoothing steps with the preconditioner. As 150 far as we know, all existing multigrid methods for fourth order problems (except 151 those in [6]) use standard smoothers and their convergence is governed by (11). 152

5 Numerical Results

The numerical experiments were performed on sienna@IMA (Intel P4, 3.4 GHz 154 CPU, 2 G memory) at the Institute for Mathematics and its Applications. In the numerical experiments we take $\sigma = 5$ and the preconditioner to be a V-cycle Poisson 156

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solve with one pre-smoothing step and one post-smoothing step. (Other multigrid 157 Poisson solves can also be used, but the V(1,1) solve appears to be the most effi-158 cient.) The contraction numbers for the *V*-cycle and *W*-cycle algorithms on the unit 159 square (with two elements in the initial mesh) are reported in Tables 1 and 2. It is 160 observed that the *V*-cycle (resp. *W*-cycle) algorithm is a contraction for $m \ge 4$ (resp. 161 $m \ge 2$).

4	5	6	7	8	9	10	11	12	13	t1.1
0.212	0.126	0.0813	0.0594	0.0442	0.0332	0.0252	0.0192	0.0147	0.0114	t1.2
0.329	0.223	0.190	0.164	0.142	0.124	0.109	0.0967	0.0861	0.0771	t1.3
0.412	0.342	0.308	0.279	0.255	0.234	0.217	0.203	0.190	0.179	t1.4
0.479	0.420	0.386	0.357	0.334	0.314	0.296	0.282	0.266	0.257	t1.5
0.537	0.467	0.434	0.408	0.386	0.367	0.351	0.336	0.324	0.312	t1.6
0.578	0.494	0.462	0.436	0.415	0.396	0.380	0.366	0.353	0.341	t1.7
0.619	0.503	0.472	0.446	0.425	0.406	0.391	0.376	0.364	0.351	t1.8
	0.212 0.329 0.412 0.479 0.537 0.578	0.212 0.126 0.329 0.223 0.412 0.342 0.479 0.420 0.537 0.467 0.578 0.494	0.212 0.126 0.0813 0.329 0.223 0.190 0.412 0.342 0.308 0.479 0.420 0.386 0.537 0.467 0.434 0.578 0.494 0.462	0.212 0.126 0.0813 0.0594 0.329 0.223 0.190 0.164 0.412 0.342 0.308 0.279 0.479 0.420 0.386 0.357 0.537 0.467 0.434 0.408 0.578 0.494 0.462 0.436	0.212 0.126 0.0813 0.0594 0.0442 0.329 0.223 0.190 0.164 0.142 0.412 0.342 0.308 0.279 0.255 0.479 0.420 0.386 0.357 0.334 0.537 0.467 0.434 0.408 0.386 0.578 0.494 0.462 0.436 0.415	0.212 0.126 0.0813 0.0594 0.0442 0.0332 0.329 0.223 0.190 0.164 0.142 0.124 0.412 0.342 0.308 0.279 0.255 0.234 0.479 0.420 0.386 0.357 0.334 0.314 0.537 0.467 0.434 0.408 0.386 0.367 0.578 0.494 0.462 0.436 0.415 0.396	0.212 0.126 0.0813 0.0594 0.0442 0.0332 0.0252 0.329 0.223 0.190 0.164 0.142 0.124 0.109 0.412 0.342 0.308 0.279 0.255 0.234 0.217 0.479 0.420 0.386 0.357 0.334 0.314 0.296 0.537 0.467 0.434 0.408 0.386 0.367 0.351 0.578 0.494 0.462 0.436 0.415 0.396 0.380	0.212 0.126 0.0813 0.0594 0.0442 0.0332 0.0252 0.0192 0.329 0.223 0.190 0.164 0.142 0.124 0.109 0.0967 0.412 0.342 0.308 0.279 0.255 0.234 0.217 0.203 0.479 0.420 0.386 0.357 0.334 0.314 0.296 0.282 0.537 0.467 0.434 0.408 0.386 0.367 0.351 0.336 0.578 0.494 0.462 0.436 0.415 0.396 0.380 0.366	0.212 0.126 0.0813 0.0594 0.0442 0.0332 0.0252 0.0192 0.0147 0.329 0.223 0.190 0.164 0.142 0.124 0.109 0.0967 0.0861 0.412 0.342 0.308 0.279 0.255 0.234 0.217 0.203 0.190 0.479 0.420 0.386 0.357 0.334 0.314 0.296 0.282 0.266 0.537 0.467 0.434 0.408 0.386 0.367 0.351 0.336 0.324 0.578 0.494 0.462 0.436 0.415 0.396 0.380 0.366 0.353	0.212 0.126 0.0813 0.0594 0.0442 0.0332 0.0252 0.0192 0.0147 0.0114 0.329 0.223 0.190 0.164 0.142 0.124 0.109 0.0967 0.0861 0.0771 0.412 0.342 0.308 0.279 0.255 0.234 0.217 0.203 0.190 0.179 0.479 0.420 0.386 0.357 0.334 0.314 0.296 0.282 0.266 0.257 0.537 0.467 0.434 0.408 0.386 0.367 0.351 0.336 0.324 0.312

Table 1. Contraction numbers for the V-cycle algorithm on the unit square.

Table 2. Contraction numbers for the W-cycle algorithm on the unit square.

k m	2	3	4	5	6	7	8	9	10	11	t2.1
1	0.661	0.368	0.212	0.126	0.0813	0.0594	0.0442	0.0332	0.0252	0.0192	t2.2
2	0.483	0.360	0.291	0.241	0.203	0.172	0.148	0.128	0.112	0.0983	t2.3
3	0.475	0.375	0.335	0.282	0.263	0.229	0.215	0.195	0.182	0.171	t2.4
4	0.455	0.383	0.335	0.308	0.287	0.270	0.256	0.244	0.233	0.223	t2.5
5	0.456	0.384	0.344	0.315	0.297	0.279	0.267	0.255	0.245	0.237	t2.6
6	0.455	0.384	0.344	0.316	0.297	0.280	0.268	0.256	0.248	0.239	t2.7
7	0.455	0.384	0.344	0.317	0.297	0.281	0.269	0.258	0.248	0.240	t2.8

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For comparison we report in Table 3 the contraction numbers for the V-cycle 163 algorithm that does not use a preconditioner in the smoothing steps. The smoothing 164 step in this algorithm is the standard Richardson relaxation scheme. 165

We have also carried out numerical experiments for the *L*-shaped domain with the vertices (0,0), (1,0), (1,1), (-1,1), (-1,-1) and (0,-1). The initial mesh consists the of six isosceles triangles sharing (0,0) as a common vertex. The contraction numbers the W-cycle algorithm with/without the preconditioner are presented in Tables 4 the and 5.

We note that the contraction numbers in Table 1 (resp. Table 4) for m smoothing 171 steps are comparable to the contraction numbers in Table 3 (resp. Tables 5) for m^2 172 smoothing steps. 173

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	m k	21	22	23	24	25	26	27	28	29	30		t3.1
	1	0.428	0.410	0.392	0.376	0.361	0.346	0.332	0.320	0.307	0.296		t3.2
	2	0.646	0.614	0.583	0.555	0.529	0.504	0.481	0.459	0.439	0.420		t3.3
	3	0.770	0.728	0.690	0.654	0.621	0.591	0.562	0.535	0.510	0.487	X	t3.4
	4	0.844	0.797	0.753	0.713	0.676	0.641	0.609	0.579	0.551	0.525		t3.5
	5	0.895	0.843	0.795	0.752	0.711	0.674	0.639	0.607	0.577	0.548		t3.6
	6	0.931	0.876	0.826	0.780	0.737	0.697	0.661	0.627	0.595	0.565		t3.7

t3.8

 Table 3. Contraction numbers for the V-cycle algorithm without a preconditioner on the unit square.

Table 4. Contraction numbers for the W-cycle algorithm with a preconditioner on the L-shape	d
domain.	

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0.960 0.902 0.849 0.801 0.757 0.715 0.677 0.642 0.609 0.578

k m	3	5	7	9	11	13	15	17	19	21	23	t4.1
1	0.319	0.187	0.125	0.105	0.0913	0.0798	0.0699	0.0614	0.0540	0.0476	0.0420	t4.2
2	0.383	0.273	0.206	0.161	0.139	0.132	0.125	0.119	0.113	0.108	0.103	t4.3
3	0.390	0.302	0.238	0.208	0.182	0.163	0.152	0.148	0.144	0.141	0.137	t4.4
4	0.386	0.309	0.271	0.245	0.224	0.208	0.193	0.181	0.170	0.161	0.153	t4.5
5	0.384	0.315	0.279	0.255	0.237	0.222	0.209	0.198	0.189	0.180	0.172	t4.6
6	0.384	0.316	0.281	0.257	0.240	0.226	0.213	0.203	0.193	0.185	0.177	t4.7
7	0.387	0.317	0.281	0.258	0.240	0.226	0.214	0.203	0.194	0.186	0.178	t4.8

 Table 5. Contraction numbers for the W-cycle algorithm without a preconditioner on the L-shaped domain.

	m k	5	7	9	11	13	15	17	19	21	23	t5.1
				0.680								t5.2
5	2	0.790	0.585	0.505	0.459	0.426	0.394	0.375	0.358	0.342	0.328	t5.3
	3	0.666	0.512	0.469	0.456	0.434	0.416	0.400	0.386	0.373	0.362	t5.4
	4	0.580	0.519	0.484	0.454	0.434	0.418	0.405	0.394	0.385	0.376	t5.5
	5	0.581	0.527	0.491	0.465	0.444	0.427	0.414	0.402	0.392	0.384	t5.6
	6	0.587	0.531	0.494	0.467	0.446	0.429	0.415	0.404	0.394	0.386	t5.7
	7	0.587	0.530	0.493	0.467	0.446	0.429	0.415	0.404	0.394	0.386	t5.8

Finally we compare the computational cost between the preconditioned schemes 174 and the un-preconditioned schemes. On the unit square, the contraction numbers for 175 the preconditioned V-cycle algorithm with m = 4 (cf. Table 1) are about the same as 176 the contraction numbers for the un-preconditioned V-cycle algorithm with m = 29 177 (cf. Table 3). For k = 7, the former takes 1.4×10^8 floating point operations and 178 0.55 s while the latter takes 3.2×10^8 floating point operations and 1.2 s. 179 Susanne C. Brenner, Shiyuan Gu, and Li-yeng Sung

On the L-shaped domain, the contraction numbers for the preconditioned W-cycle 180 algorithm with m = 3 (cf. Table 4) are about the same as the contraction numbers 181 for the un-preconditioned W-cycle algorithm with m = 23 (cf. Table 5). For k = 7, 182 the former takes 4.7×10^8 floating point operations and 2.1 s while the latter takes 183 1.1×10^9 floating point operations and 4.7 s. 184

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