
A Substructuring Preconditioner for Three-Dimensional Maxwell's Equations

Qiya Hu¹, Shi Shu² and Jun Zou³

¹ LSEC, Institute of Computational Mathematics and Scientific/Engineering Computing, Academy of Mathematics and Systems Science, The Chinese Academy of Sciences, Beijing 100080, China hqy@lsec.cc.ac.cn

² School of Mathematics and Computational Science, Xiangtan University, Hunan, 411105, China shushi@xtu.edu.cn

³ Department of Mathematics, The Chinese University of Hong Kong, Shatin, N.T., Hong Kong zou@math.cuhk.edu.hk

Summary. We propose a new nonoverlapping domain decomposition preconditioner for the discrete system arising from the edge element discretization of the three-dimensional Maxwell's equations. This preconditioner uses the simplest coarse edge element space induced by the coarse triangulation. We will show that the rate of the PCG convergence with this substructuring preconditioner is quasi-optimal, and is independent of large variations of the coefficients across the local interfaces.

1 Introduction

When the time-dependent Maxwell's equations is solved numerically, we need to solve the following **curlcurl**-system at each time step [4, 6, 8, 12]:

$$\operatorname{curl}(\alpha \operatorname{curl} \mathbf{u}) + \beta \mathbf{u} = \mathbf{f} \quad \text{in } \Omega \quad (1)$$

where Ω is assumed to be an open polyhedral domain in \mathbf{R}^3 , and the coefficients $\alpha(\mathbf{x})$ and $\beta(\mathbf{x})$ are two positive bounded functions in Ω . We shall complement the Eq. (1) with the perfect conductor condition $\mathbf{u} \times \mathbf{n} = 0$ on $\partial\Omega$, where \mathbf{n} is the unit outward normal vector on $\partial\Omega$.

Edge finite element methods have been widely applied in the numerical solution of the system (1), see, for example, [5, 6, 8, 11]. Compared to the standard nodal finite element methods, the discrete systems resulting from the edge element discretization are essentially different in nature. The non-overlapping domain decomposition preconditioners have been well developed for the nodal element systems for the standard second order elliptic problems in the past two decades, and proved both numerically and theoretically to perform nearly optimally in terms of the fine mesh size and subdomain size; see, e.g., the monograph [15]. But these preconditioners,

or their natural generalizations turn out to perform mostly very poorly for the edge 33
element systems for the **curlcurl**-system (1), especially in three dimensions. 34

A lot of important efforts have been made in the construction of effective do- 35
main decomposition methods for the system (1). A substructuring type method was 36
analysed in [16] for two dimensions, and in [2] for three dimensions with two sub- 37
domains. In [7], a novel substructuring type method was proposed for general two- 38
dimensional multiple subdomains with quite irregular boundaries, and it was proved 39
to be nearly optimal in terms of a variety of mesh decompositions and distributions 40
of physical material properties. However, it has been a challenge how to construct a 41
efficient non-overlapping domain decomposition preconditioner for the Maxwell's 42
equations in three dimensions with general multiple subdomains. A first important 43
attempt to this problem was made in [9] where a wire basket type algorithm was pro- 44
posed and analysed. Then a substructuring preconditioner and a dual-primal FETI 45
algorithm were introduced and fully analysed for three dimensions in [10] and [14], 46
respectively. These three methods have their respective advantages and disadvan- 47
tages: the algorithms in [9] and [14] both involve smaller coarse solvers but they are 48
difficult to implement; the method in [10] is easier to implement but it involves a 49
relatively large coarse solver. 50

This work intends to construct a new substructuring type preconditioner for the 51
three-dimensiona**l curlcurl**-system (1) for general multiple subdomains. In this pre- 52
conditioner, the coarse space is chosen to be the edge element space induced by 53
the coarse triangulation, so the resulting coarse solver is very cheap and simple to 54
implement. It is shown that the rate of the PCG convergence with this substructur- 55
ing preconditioner is quasi-optimal, and more importantly, independent of the large 56
variations of the coefficients in the system (1) across the local interfaces. 57

2 Domain Decompositions and Discretizations 58

This section introduces the non-overlapping domain decomposition of domain Ω , 59
the weak form of the system (1) and the edge element spaces. 60

2.1 Initial Domain Decomposition Based on the Distribution of the Coefficients 61

We assume that the entire domain Ω is decomposed into N_0 open convex polyhedral 62
subdomains D_1, D_2, \dots, D_{N_0} such that $\bar{\Omega} = \cup_{r=1}^{N_0} \bar{D}_r$ and $\alpha(\mathbf{x})$ and $\beta(\mathbf{x})$ are positive 63
constants on each subdomain D_r , namely for $r = 1, 2, \dots, N_0$, 64

$$\alpha(\mathbf{x}) = \alpha_r, \quad \beta(\mathbf{x}) = \beta_r \quad \forall \mathbf{x} \in D_r. \quad 65$$

Clearly such a decomposition is always possible when the domain Ω is occupied by 66
multiple media. In fact, if for some medium we have an irregular nonconvex subre- 67
gion in Ω , we can further split each nonconvex medium subregion into smaller con- 68
vex subdomains. This means that our assumption does cover many practical cases, 69
especially considering the fact that the domain Ω on which we solve the original 70

Maxwell system (1) by a finite element method is often obtained by approximating the original physical domain by a polyhedral domain. Note that N_0 typically is a fixed constant in applications, so $\text{diam}(D_r) = O(1)$.

Let F_{nm} denote the common face of two neighboring subdomains D_n and D_m , and set $D_{nm} = D_n \cup D_m \cup F_{nm}$. For simplicity of the analysis, we assume

$$\beta_r \lesssim \alpha_r \lesssim d^{-2} \alpha_r, \quad r = 1, \dots, N_0. \quad (2)$$

2.2 Domain Decomposition

For a number $d \in (0, 1)$, let each polyhedron D_l be decomposed into the union of some non-overlapping tetrahedra (or hexahedra) $\{\Omega_k\}$ of size d (see [3, 15] and [18]), which results in a non-overlapping domain decomposition for Ω : $\bar{\Omega} = \bigcup_{k=1}^N \bar{\Omega}_k$.

Naturally we further assume that $\Omega_i \cap \Omega_j = \emptyset$ when $i \neq j$; if $i \neq j$ and $\partial\Omega_i \cap \partial\Omega_j \neq \emptyset$, $\partial\Omega_i \cap \partial\Omega_j$ is a common face (or edge or vertex) of Ω_i and Ω_j . Now the subdomains $\Omega_1, \dots, \Omega_N$ constitute our desired coarse triangulation \mathcal{T}_d of Ω . The faces and vertices of the subdomains are always denoted by F and v , while the common (open) face of the subdomains Ω_i and Ω_j are denoted by Γ_{ij} , and the union of all such common faces by Γ , i.e., $\Gamma = \cup \Gamma_{ij}$. Γ will be called the interface. By Γ_k we denote the intersection of Γ with the boundary of the subdomain Ω_k . So we have $\Gamma_k = \partial\Omega_k$ if Ω_k is an interior subdomain of Ω . We shall set $\Omega_{ij} = \Omega_i \cup \Omega_j \cup \Gamma_{ij}$.

2.3 Weak Formulation

Let $H(\mathbf{curl}; \Omega)$ be the Sobolev space consisting of all square integrable functions whose \mathbf{curl} 's are also square integrable in Ω , and $H_0(\mathbf{curl}; \Omega)$ be a subspace of $H(\mathbf{curl}; \Omega)$ of all functions whose tangential components vanish on $\partial\Omega$. Then by writing the scalar product in $(L^2(\Omega))^3$ as (\cdot, \cdot) , we can state the variational problem for system (1) as follows:

Find $\mathbf{u} \in H_0(\mathbf{curl}; \Omega)$ such that

$$\mathcal{A}(\mathbf{u}, \mathbf{v}) = (\mathbf{f}, \mathbf{v}), \quad \forall \mathbf{v} \in H_0(\mathbf{curl}; \Omega) \quad (3)$$

where $\mathcal{A}(\cdot, \cdot)$ is a bilinear form given by

$$\mathcal{A}(\mathbf{u}, \mathbf{v}) = (\alpha \mathbf{curl} \mathbf{u}, \mathbf{curl} \mathbf{v}) + (\beta \mathbf{u}, \mathbf{v}), \quad \mathbf{u}, \mathbf{v} \in H(\mathbf{curl}; \Omega).$$

2.4 Fine Triangulation and Their Associated Finite Element Spaces

We further divide each Ω_k into smaller tetrahedral elements of size h so that elements from two neighboring subdomains have an intersection which is either empty or a single nodal point or an edge or a face on the interface Γ . Let \mathcal{T}_h be the resulting triangulation of the domain Ω , which we assume is quasi-uniform. Then we introduce the Nédélec edge element space of the lowest order defined on \mathcal{T}_h (cf. [12] and [13]):

$$V_h(\Omega) = \left\{ \mathbf{v} \in H_0(\mathbf{curl}; \Omega); \mathbf{v}|_K \in R(K), \forall K \in \mathcal{T}_h \right\}, \quad 104$$

where $R(K)$ is a subset of all linear polynomials on the element K of the form: 105

$$R(K) = \left\{ \mathbf{a} + \mathbf{b} \times \mathbf{x}; \mathbf{a}, \mathbf{b} \in \mathbf{R}^3, \mathbf{x} \in K \right\}. \quad 106$$

In an analogous way, we can define the coarse edge element space $V_d(\Omega) \subset V_h(\Omega)$, 107
associated with the *coarse* triangulation \mathcal{T}_d . 108

It is well-known that for any $\mathbf{v} \in V_h(\Omega)$, its tangential components are continuous 109
on all edges of each element in the triangulation \mathcal{T}_h . Moreover, each edge element 110
function \mathbf{v} in $V_h(\Omega)$ is uniquely determined by its moments on each edge e of \mathcal{T}_h : 111

$$\left\{ \lambda_e(\mathbf{v}) = \int_e \mathbf{v} \cdot \mathbf{t}_e ds; e \in \mathcal{E}_h \right\}, \quad 112$$

where \mathcal{E}_h denotes the set of the *fine* edges from the triangulation \mathcal{T}_h , and \mathbf{t}_e denotes 113
the unit vector on the edge e . 114

By $Z_h(\Omega)$ we denote the continuous piecewise linear finite element subspace of 115
 $H_0^1(\Omega)$ associated with the triangulation \mathcal{T}_h . Similarly, let $Z_d(\Omega)$ denote the contin- 116
uous piecewise linear finite element subspace of $H_0^1(\Omega)$ associated with the triangu- 117
lation \mathcal{T}_d . 118

2.5 Discrete Variational Problem 119

Using the edge element space $V_h(\Omega)$, the system (3) may be approximated as fol- 120
lows: Find $\mathbf{u}_h \in V_h(\Omega)$ such that 121

$$(\alpha \mathbf{curl} \mathbf{u}_h, \mathbf{curl} \mathbf{v}_h) + (\beta \mathbf{u}_h, \mathbf{v}_h) = (\mathbf{f}, \mathbf{v}_h), \quad \forall \mathbf{v}_h \in V_h(\Omega). \quad (4)$$

Define the operator $A : V_h(\Omega) \rightarrow V_h(\Omega)$ by 122

$$(A\mathbf{u}_h, \mathbf{v}_h) = (\alpha \mathbf{curl} \mathbf{u}_h, \mathbf{curl} \mathbf{v}_h) + (\beta \mathbf{u}_h, \mathbf{v}_h), \quad \forall \mathbf{u}_h, \mathbf{v}_h \in V_h(\Omega), \quad 123$$

Then, (4) can be written in the operator form 124

$$A\mathbf{u}_h = \mathbf{f}_h. \quad (5)$$

3 A Nearly Optimal Preconditioner for A 125

3.1 Construction of the Preconditioner 126

We first introduce some useful sets and subspaces. 127

\mathcal{E}_h : the set of all edges from the triangulations \mathcal{T}_h ; 128

$\mathcal{E}_{\Gamma,h}$: the set of edges which belong to \mathcal{E}_h and have two endpoints on the interface 129

Γ ; 130

\mathcal{E}_d : the set of all (coarse) edges from the triangulations \mathcal{T}_d ; 131

\mathscr{W}_E : the union of all the coarse edges $E' \in \mathcal{E}_d$, which have a common endpoint with the coarse edge $E \in \mathcal{E}_d$. And \mathscr{W}_E is called *E-basket*.

$\mathcal{E}_{E,h}^b$: the set of all (fine) edges which belong to \mathcal{E}_h and have at least one endpoint on \mathscr{W}_E ;

Let D be either a subdomain D_r or a subdomain Ω_k or a subdomain Ω_{ij} or a subdomain D_{mn} . The restrictions of $V_h(\Omega)$ (resp. $Z_h(\Omega)$) on D is denoted by $V_h(D)$ (resp. $Z_h(D)$). The following local subspaces of $V_h(D)$ will be important to our analysis:

$$V_h^0(D) = \left\{ \mathbf{v} \in V_h(D); \mathbf{v} \times \mathbf{n} = 0 \text{ on } \partial D \right\},$$

and

$$Z_h^0(D) = \left\{ \varphi \in Z_h(\Omega); \text{supp } \varphi \subset D \right\}.$$

We define subspaces of $V_h(\Omega)$:

$$V_h^H(\Omega) = \left\{ \mathbf{v} \in V_h(\Omega); \mathbf{v} \text{ is the discrete } A\text{-extension of } \mathbf{v}|_{\partial\Omega_k} \text{ in each } \Omega_k \right\},$$

$$V_h^H(\Omega_{ij}) = V_h^H(\Omega) \cap V_h^0(\Omega_{ij}),$$

and for $E \in \mathcal{E}_d$,

$$V_h^E(\Omega) = \left\{ \mathbf{v} \in V_h^H(\Omega); \lambda_e(\mathbf{v}) = 0 \text{ for each } e \in \mathcal{E}_{\Gamma,h} \setminus \mathcal{E}_{E,h}^b \right\}.$$

It is well known that a suitable *coarse* subspace plays a key role in the construction of an effective domain decomposition preconditioner, and it is generally rather technical and problem-dependent to choose such a *coarse* subspace. Surprisingly we are going to choose the coarse subspace to be the simplest one, namely the subspace $V_d(\Omega)$ induced by the coarse triangulation \mathcal{T}_d .

It is easy to see that the space $V_h(\Omega)$ has the (non-direct sum) decomposition

$$V_h(\Omega) = V_d(\Omega) + \sum_{k=1}^N V_h^0(\Omega_k) + \sum_E V_h^E(\Omega) + \sum_{\Gamma_{ij}} V_h^H(\Omega_{ij}). \quad (6)$$

Next, we define the corresponding solvers on the subspaces $V_h^0(\Omega_k)$, $V_h^E(\Omega)$, $V_h^H(\Omega_{ij})$ and $V_d(\Omega)$.

As usual, we denote the restriction of A on $V_h^0(\Omega_k)$ by A_k , i.e.,

$$(A_k \mathbf{v}, \mathbf{u})_{\Omega_k} = (A \mathbf{v}, \mathbf{u}) = \mathcal{A}(\mathbf{v}, \mathbf{u}), \quad \mathbf{v} \in V_h^0(\Omega_k), \quad \forall \mathbf{u} \in V_h^0(\Omega_k).$$

Let $B_k : V_h^0(\Omega_k) \rightarrow V_h^0(\Omega_k)$, $B_d : V_d(\Omega) \rightarrow V_d(\Omega)$ and $B_{ij} : V_h^H(\Omega_{ij}) \rightarrow V_h^H(\Omega_{ij})$ be the symmetric and positive definite operators such that

$$(B_k \mathbf{v}, \mathbf{v}) \cong (A_k \mathbf{v}_k, \mathbf{v}_k)_{\Omega_k}, \quad \forall \mathbf{v} \in V_h^0(\Omega_k),$$

where $\mathbf{v}_k = \mathbf{v}|_{\Omega_k}$ for $k = 1, 2, \dots, N$, and

$$\begin{aligned} (B_d \mathbf{v}_d, \mathbf{v}_d) &\cong \mathcal{A}(\mathbf{v}_d, \mathbf{v}_d), \quad \forall \mathbf{v}_d \in V_d(\Omega), \\ (B_{ij} \mathbf{v}, \mathbf{v}) &\cong \mathcal{A}(\mathbf{v}, \mathbf{v}), \quad \forall \mathbf{v} \in V_h^H(\Omega_{ij}). \end{aligned}$$

The symbol \cong above means each of the two quantities involved is bounded by the other up to a constant independent of h, d and functions involved in the two quantities.

The local solvers on $V_h^E(\Omega)$ should be solvable in an efficient manner, and their constructions are much more tricky and technical than the others. To do so, we introduce more notation.

For any face F from the triangulations \mathcal{T}_d , we use F_b to denote the union of all \mathcal{T}_h -induced (closed) triangles on F , which have either one single vertex or one edge lying on ∂F , and F_∂ to denote the open set $F \setminus F_b$. For any subdomain Ω_k , define

$$\Delta_k = \bigcup_{F \subset \Gamma_k} F_b, \quad k = 1, \dots, N. \quad (171)$$

We will also need the so-called tangential divergence $\text{div}_\tau \Phi = \text{curl}_S \Phi$ for $\Phi \in V_h(\Gamma_k)$, which is defined here as in [1, 2]. Then we can introduce our local solver $B_E : V_h^E(\Omega) \rightarrow V_h^E(\Omega)$ as follows:

$$\begin{aligned} (B_E \mathbf{v}, \mathbf{u}) &= h[1 + \log(d/h)] \sum_{k=1}^N \left\{ \alpha_k \langle \text{div}_\tau(\mathbf{v} \times \mathbf{n})|_{\Gamma_k}, \text{div}_\tau(\mathbf{u} \times \mathbf{n})|_{\Gamma_k} \rangle_{\Delta_k} \right. \\ &\quad \left. + \beta_k \langle \mathbf{v} \times \mathbf{n}, \mathbf{u} \times \mathbf{n} \rangle_{\Delta_k} \right\}, \quad \mathbf{v} \in V_h^E(\Omega), \quad \forall \mathbf{u} \in V_h^E(\Omega). \end{aligned} \quad (7)$$

For convenience, we call B_E an

Let $Q_k : V_h(\Omega) \rightarrow V_h^0(\Omega_k)$, $Q_d : V_h(\Omega) \rightarrow V_d(\Omega)$, $Q_E : V_h(\Omega) \rightarrow V_h^E(\Omega)$ and $Q_{ij} : V_h(\Omega) \rightarrow V_h^H(\Omega_{ij})$ be the standard the standard L^2 -projections. Then we are ready to propose our new preconditioner for A as follows:

$$B^{-1} = B_d^{-1} Q_d + \sum_{k=1}^N B_k^{-1} Q_k + \omega \sum_E B_E^{-1} Q_E + \sum_{I_{ij}} B_{ij}^{-1} Q_{ij}, \quad (8)$$

where ω is a (constant) relaxation parameter, which is introduced to obtain a balance between the local solvers B_E and other remaining solvers.

3.2 Algorithm Based on the New Preconditioner and Main Results

The action of the preconditioner B^{-1} which is needed in each PCG iteration can be described in the following algorithm.

Algorithm 4.1. For $\mathbf{g} \in V_h(\Omega)$, we can compute $\mathbf{u} = B^{-1} \mathbf{g}$ in five steps.

Step 1. Solve the system for $\mathbf{u}_d \in V_d(\Omega)$:

$$(B_d \mathbf{u}_d, \mathbf{v}_d) = (\mathbf{g}, \mathbf{v}_d), \quad \forall \mathbf{v}_d \in V_d(\Omega);$$

Step 2. Solve the following system for $\mathbf{u}_k \in V_h^0(\Omega_k)$ in each subdomain in parallel: 187
188

$$(B_k \mathbf{u}_k, \mathbf{v}) = (\mathbf{g}, \mathbf{v}), \quad \forall \mathbf{v} \in V_h^0(\Omega_k), \quad k = 1, \dots, N; \quad 189$$

Step 3. Solve the following system for $\mathbf{u}_{ij} \in V_h^0(\Omega_{ij})$ in each subdomain Ω_{ij} in parallel: 190
191

$$(B_{ij} \mathbf{u}_{ij}, \mathbf{v}) = (\mathbf{g}, \mathbf{v}) - (A_i \mathbf{u}_i, \mathbf{v})_{\Omega_i} - (A_j \mathbf{u}_j, \mathbf{v})_{\Omega_j}, \quad \forall \mathbf{v} \in V_h^0(\Omega_{ij}); \quad 192$$

Step 4. Solve the system for $\mathbf{u}_E \in V_h^E(\Omega)$: 193

$$(B_E \mathbf{u}_E, \mathbf{v}) = (\mathbf{g}, \tilde{\mathbf{v}}) - \sum_{k=1}^N (A_k \mathbf{u}_k, \tilde{\mathbf{v}}), \quad \mathbf{v} \in V_h^E(\Omega), \quad 194$$

where $\tilde{\mathbf{v}} \in V_h(\Omega)$ is a natural extension of $(\mathbf{v} \times \mathbf{n})|_\Gamma$ by zero. 195

Step 5. Set $\Phi_h = (\sum_{ij} \mathbf{u}_{ij} + \sum_E \mathbf{u}_E) \times \mathbf{n}|_\Gamma$ and compute the A -extension 196

of Φ_h on each Ω_k to obtain $\mathbf{u}^H \in V_h^H(\Omega)$. This leads to 197

$$\mathbf{u} = \mathbf{u}_d + \sum_{k=1}^N \mathbf{u}_k + \mathbf{u}^H. \quad 198$$

Remark 1. For the local solver B_{ij} on each face Γ_{ij} , we may use the face extended 199
domain formed by, e.g., one half of each of the two neighboring subdomains Ω_i 200
and Ω_j . Such definition of B_{ij} 's can reduce the computational complexity in their 201
numerical realization. 202

Let E denote a coarse edge of the subdomain D_r . Define 203

$$V_h^\perp(\Omega) = \{\mathbf{v}_h : \mathbf{v}_h \in V_h(\Omega), \int_E \mathbf{v}_h \cdot \mathbf{t}_E ds = 0 \text{ for each } E\}. \quad 204$$

We shall use $\kappa^\perp(B^{-1}A)$ to denote the *induced condition number* of the preconditioned system $B^{-1}A$ associated with the subspace $V_h^\perp(\Omega)$, namely the condition number of $B^{-1}A$ restricted on the subspace $V_h^\perp(\Omega)$ (cf. [17]). At this moment we are able to establish only the following estimate of the induced condition number. As the estimate is quite lengthy and technical, we cannot include it here due to the page limitation. 205
206
207
208
209
210

Theorem 1. *Under the assumptions (2), the preconditioner B given in (8) is nearly optimal in the sense that* 211
212

$$\kappa^\perp(B^{-1}A) \leq C[1 + \log(d/h)]^2[1 + \log(1/h)]^2 \quad (9)$$

where the constant C is independent of h , d and the jumps of the coefficients. 213

As we see from the above theorem that the induced condition number grows logarithmically with the degrees of freedom in each subdomain, but also with the 214
215

degrees of freedom of the entire fine mesh. We believe this is mainly due to the re- 216
restriction of our current analysis technique, namely the estimate must be done for the 217
induced condition number in the subspace $V_h^\perp(\Omega)$ associated with the coarse trian- 218
gulation formed by the material subdomains D_r . We expect the estimate should be 219
finally carried out directly in the entire edge element space $V_h(\Omega)$, that will remove 220
the logarithmic factor of $1/h$ in the estimate (9). This expectation has already been 221
confirmed by our three-dimensional numerical experiments; see the next section. 222

4 Numerical Experiments 223

In this section we shall conduct some numerical experiments to check the conver- 224
gence of the newly proposed preconditioner, and find out whether they are consistent 225
with the prediction of the convergence theory developed in the previous sections. 226

In our experiments, we take the domain to be the unit cube $\Omega = (0, 1)^3$, while 227
the right-hand side \mathbf{f} of the system (1) is selected such that the exact solution $\mathbf{u} =$ 228
 $(u_1, u_2, u_3)^T$ is given by 229

$$\begin{aligned} u_1 &= xyz(x-1)(y-1)(z-1), \\ u_2 &= \sin(\pi x) \sin(\pi y) \sin(\pi z), \\ u_3 &= (1-e^x)(1-e^{x-1})(1-e^y)(1-e^{y-1})(1-e^z)(1-e^{z-1}), \end{aligned}$$

when the coefficients $\alpha(\mathbf{x})$ and $\beta(\mathbf{x})$ are both constant 1. This right-hand side \mathbf{f} is 230
then fixed in all our experiments, but the coefficients $\alpha(\mathbf{x})$ and $\beta(\mathbf{x})$ may be taken 231
differently. 232

We then need to triangulate the domain Ω into subdomains $\{\Omega_k\}$. For this, we 233
first partition the three edges of Ω on x -, y - and z -axis into n equal subintervals from 234
which one can naturally generate n^3 equal smaller cubes of size $d = 1/n$. This yields 235
the desired subdomain decomposition in our experiments. 236

Next, we further triangulate each subdomain Ω_k to get a fine triangulation \mathcal{T}_h of 237
size h over the domain Ω . To generate \mathcal{T}_h , we divide each subdomain into m^3 equal 238
smaller cubes of size $h = 1/(mn)$, in the same manner as done in the previous subdo- 239
main generation. Then \mathcal{T}_h is obtained by triangulating each cube into six tetrahedra. 240
For easy identification, we may denote the triangulation \mathcal{T}_h as $m^3(n^3)$ below. 241

The edge finite element space of the lowest order is used for the discretization 242
of (3). The resulting system (5) is solved by PCG method with the newly proposed 243
preconditioners B defined in Sect. 4. We shall choose the balancing parameter ω in 244
front of the *E-basket* local solvers B_E in (8) as $\omega = 1$ or $\omega = 2.5$. 245

We consider various distributions of the coefficients $\alpha(\mathbf{x})$ and $\beta(\mathbf{x})$ and report 246
the corresponding numbers of PCG iterations, and the condition numbers of $B^{-1}A$ 247
for some representative cases. The PCG iteration is terminated in our experiments 248
when the relative residual is less than 10^{-6} . 249

Case (i): coefficients $\alpha(\mathbf{x}) = \beta(\mathbf{x}) = 1$, with no jumps. The PCG iterations and the 250
condition numbers (in brackets) for $\omega = 2.5$ are listed in Table 1. 251

$m \setminus n$	$\omega = 1.0$				$\omega = 2.5$			
	4	6	8	10	4	6	8	10
4	34	33	32	32	31 (34.24)	31 (36.31)	31 (36.94)	30 (37.40)
8	41	40	39	38	39 (52.15)	38 (53.78)	37 (54.21)	37 (54.61)
12	48	47	44	42	43 (64.29)	43 (65.91)	41 (66.19)	41 (66.62)
16	51	50	49	45	47 (74.40)	46 (75.69)	44 (75.82)	44 (76.39)

Table 1. Iterations (and condition numbers) with smooth coefficients

We observe from the above table that the number of PCG iterations grows slowly when $m = d/h$ increases but $n = 1/d$ is fixed, and that these numbers vary stably when m is fixed but n increases. This justifies our early expectation that the condition number of the preconditioned system $B^{-1}A$ should grow logarithmically with d/h only, not with $1/h$.

One important issue we like to draw the readers' attention to is the large-scale of the discrete system we are solving. For instance, when $m = 16$ and $n = 10$, the total number of degrees of freedom for the fine edge element system is about 28,672,000.

Case (ii): coefficients $\alpha(\mathbf{x})$ and $\beta(\mathbf{x})$ have large jumps:

$$\alpha(\mathbf{x}) = \beta(\mathbf{x}) = \alpha_0 \quad \text{in } D; \quad \alpha(\mathbf{x}) = \beta(\mathbf{x}) = 1 \quad \text{in } \Omega \setminus D.$$

where $D \subset \Omega$ is a union of several subdomains Ω_k . We choose $\alpha_0 = 10^{-5}$ or $\alpha_0 = 10^5$, and consider two choices of D , where one does not have *cross-points*, while the other has one *cross-point*.

Example 1:

$$D = \left[\frac{1}{4}, \frac{1}{2}\right]^3.$$

Example 2:

$$D = \left[\frac{1}{4}, \frac{1}{2}\right]^3 \cup \left[\frac{1}{2}, \frac{3}{4}\right]^3.$$

The numerical results are given in Tables 2 and 3, from which we can make some similar observations about the PCG convergence in terms of the mesh and subdomain quantities d/h and d as we did for Case (i).

$m \setminus n$	Example 1				Example 2			
	$\omega = 1.0$		$\omega = 2.5$		$\omega = 1.0$		$\omega = 2.5$	
	4	8	4	8	4	8	4	8
4	29	31	26 (32.00)	29 (35.97)	28	30	26 (35.51)	30 (35.97)
8	35	38	32 (44.88)	37 (52.97)	35	38	32 (45.88)	37 (52.59)
12	38	45	36 (56.02)	42 (64.96)	37	45	35 (55.66)	41 (63.81)
16	40	49	37 (64.65)	45 (74.68)	40	49	37 (65.65)	45 (74.31)

Table 2. Iterations (and condition numbers) with $\alpha_0 = 10^{-5}$

		Example 1				Example 2			
		$\omega = 1.0$		$\omega = 2.5$		$\omega = 1.0$		$\omega = 2.5$	
$m \setminus n$	4	8	4	8	4	8	4	8	
4	42	42	36 (40.47)	36 (42.71)	42	44	38 (40.55)	37 (42.72)	
8	49	48	45 (61.08)	44 (62.89)	52	51	46 (60.20)	45 (62.89)	
12	55	54	50 (74.04)	49 (76.28)	56	56	50 (76.24)	51 (76.28)	
16	59	57	54 (91.51)	52 (86.45)	59	59	53 (83.35)	54 (86.45)	

Table 3. Iterations (and condition numbers) with $\alpha_0 = 10^5$

Case (iii): coefficients $\alpha(\mathbf{x})$ and $\beta(\mathbf{x})$ have large jumps:

$$\alpha(\mathbf{x}) = \begin{cases} \alpha_0, & \text{in } D \\ 1, & \text{in } \Omega \setminus D, \end{cases} \quad \beta(\mathbf{x}) = \begin{cases} \beta_0, & \text{in } D \\ 1, & \text{in } \Omega \setminus D, \end{cases}$$

where $D \subset \Omega$ is a union of several subdomains Ω_k . We choose $\alpha_0 = 10^{-5}$ or $\alpha_0 = 10^5$, but $\beta_0 \neq \alpha_0$. We still consider two different regions D from Examples 1 and 2 in the previous Case (ii), but choose the balancing parameter ω in front of the E-basket local solvers B_E in (8) as $\omega = 2.5$.

The numerical results are given in Tables 4 and 5. Again, we can make similar observations about the PCG convergence in terms of the mesh and subdomain quantities d/h and d as we did for Case (i).

		Example 1				Example 2			
		$\beta_0 = \alpha_0 \times 10^2$		$\beta_0 = \alpha_0 \times 10^{-2}$		$\beta_0 = \alpha_0 \times 10^2$		$\beta_0 = \alpha_0 \times 10^{-2}$	
$m \setminus n$	4	8	4	8	4	8	4	8	
4	30	36	46	47	30	36	45	47	
8	39	43	56	56	39	45	56	56	
16	49	52	65	65	49	52	63	65	

Table 4. Iterations with $\alpha_0 = 10^{-5}$

		Example 1				Example 2			
		$\beta_0 = \alpha_0 \times 10^2$		$\beta_0 = \alpha_0 \times 10^{-2}$		$\beta_0 = \alpha_0 \times 10^2$		$\beta_0 = \alpha_0 \times 10^{-2}$	
$m \setminus n$	4	8	4	8	4	8	4	8	
4	31	37	38	41	31	37	39	46	
8	37	47	46	49	37	47	53	58	
16	48	56	55	57	48	56	66	73	

Table 5. Iterations with $\alpha_0 = 10^5$

We may also observe from the previous numerical experiments that appropriate choices of the parameter ω can significantly improve the efficiency of the preconditioner B . It is important to see that the choices of ω seem independent of the fine

and coarse meshsizes h and d , so we may determine ω by solving some small scale systems, e.g., a system with $m = n = 4$.

Acknowledgments QH was supported by the Major Research Plan of Natural Science Foundation of China G91130015, the Key Project of Natural Science Foundation of China G11031006 and National Basic Research Program of China G2011309702. SS was supported by NSFC Project 91130002 and 11171281, the project of Scientific Research Fund of Hunan Provincial Education Department 10C1265 and 11C1219, and the project of Xiangtan University 10XZX03. JZ was substantially supported by Hong Kong RGC grant (Project 405110) and a Direct Grant from Chinese University of Hong Kong.

Bibliography

- [1] A. Alonso and A. Valli. Some remarks on the characterization of the space of tangential traces of $H(\mathbf{curl}; \Omega)$ and the construction of an extension operator. *Manuscr. Math.*, 89:159–178, 1996.
- [2] A. Alonso and A. Valli. An optimal domain decomposition preconditioner for low-frequency time-harmonic maxwell equations. *Math. Comp.*, 68(6):607–631, 1999.
- [3] J. Bramble, J. Pasciak, and A. Schatz. The construction of preconditioner for elliptic problems by substructuring. *IV. Math. Comp.*, 53(5):1–24, 1989.
- [4] M. Cessenat. *Mathematical methods in electromagnetism*. World Scientific, River Edge, NJ, 1998.
- [5] Z. Chen, Q. Du, and J. Zou. Finite element methods with matching and non-matching meshes for maxwell equations with discontinuous coefficients. *SIAM J. Numer. Anal.*, 37:1542–1570, 1999.
- [6] P. Ciarlet, Jr., and J. Zou. Fully discrete finite element approaches for time-dependent maxwell's equations. *Numer. Math.*, 82(8):193–219, 1999.
- [7] C. R. Dohrmann and O. Widlund. An iterative substructuring algorithm for two-dimensional problems in $H(\mathbf{curl})$. Technical report, TR2010-936, Courant Institute, New York, 2010.
- [8] R. Hiptmair. Finite elements in computational electromagnetism. *Acta Numerica*, 11:237–339, 2002.
- [9] Q. Hu and J. Zou. A non-overlapping domain decomposition method for maxwell's equations in three dimensions. *SIAM J. Numer. Anal.*, 41:1682–1708, 2003.
- [10] Q. Hu and J. Zou. Substructuring preconditioners for saddle-point problems arising from maxwell's equations in three dimensions. *Math. Comput.*, 73:35–61, 2004.
- [11] P. Monk. Analysis of a finite element method for maxwell's equations. *SIAM J. Numer. Anal.*, 29:32–56, 1992.
- [12] P. Monk. *Finite Element Methods for Maxwell's Equations*. Oxford University Press, Oxford, 2003.
- [13] J. Nedelec. Mixed finite elements in R^3 . *Numer. Math.*, 35:315–341, 1980.

- [14] A. Toselli. Dual-primal FETI algorithms for edge finite-element approximations in 3d. *IMA J. Numer. Anal.*, 26:96–130, 2006. 325
326
- [15] A. Toselli and O. Widlund. *Domain Decomposition Methods – Algorithms and Theory*. Springer, New York, 2004. 327
328
- [16] A. Toselli, O. Widlund, and B. Wohlmuth. An iterative substructuring method for Maxwell’s equations in two dimensions. *Math. Comp.*, 70:935–949, 2001. 329
330
- [17] J. Xu and Y. Zhu. Uniform convergent multigrid methods for elliptic problems with strongly discontinuous coefficients. *M³AS*, 18:77–105, 2008. 331
332
- [18] J. Xu and J. Zou. Some non-overlapping domain decomposition methods. *SIAM Review*, 40:857–914, 1998. 333
334

UNCORRECTED PROOF