

Towards a IETI-DP Solver on Non-Matching Multi-Patch Domains

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1 Introduction

Isogeometric Analysis (IgA), see [7], is a method for discretizing partial differential equations (PDEs). The goal of its development has been to enhance the interface between computer-aided design (CAD) and simulation. Current state-of-the-art CAD tools use B-splines and NURBS for the representation of the computational domain. In IgA, the same kind of bases is also utilized to discretize the PDEs. Complex domains for real-world applications are usually the union of many patches, parametrized with individual geometry functions (multi-patch IgA). We focus on non-overlapping patches.

If the grids are not conforming and/or the interfaces between the patches do not consist of whole edges then discontinuous Galerkin (dG) methods are the discretization techniques of choice. A well studied representative is the symmetric interior discontinuous Galerkin (SIPG) method, cf. [1]. It has already been adapted and analyzed in IgA, cf. [9, 10, 15] and others. An obvious choice to solve discretized PDEs on domains with many non-overlapping patches are tearing and interconnecting methods. The variant we are interested in is the dual-primal approach, see [4] for FETI-DP and [8, 5, 6] for its extension to IgA, which is called accordingly dual-primal isogeometric tearing and interconnecting method (IETI-DP). This method is similar to Balancing Domain Decomposition with Constraints (BDDC) methods, that have also been adapted to IgA, see [2, 17] and references therein. In [14, 15], the authors have presented a p - and h -robust convergence analysis for IETI-DP. The

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authors have assumed that the interfaces consist of whole edges. If the vertices are chosen as primal degrees of freedom, it was shown that the condition number of the preconditioned Schur complement system is, under proper assumptions, bounded by

$$C p \left(1 + \log p + \max_{k=1, \dots, K} \log \frac{H_k}{h_k} \right)^2, \quad (1)$$

where p is the spline degree, h_k is the grid size on patch $\Omega^{(k)}$ and H_k is the diameter of $\Omega^{(k)}$ and $C > 0$ is a constant independent of these quantities. In this paper, we construct a new IETI-DP method that can deal with interfaces that do not consist of whole edges. This means that the patches can meet in T-junctions, which increases the flexibility of the geometric model significantly. In this IETI-DP variant, the construction of the coarse space is based on the idea of ‘‘fat vertices’’: We consider every basis function that is supported on a vertex or T-junction as primal degree of freedom. The numerical experiments indicate that a similar condition number bound to (1) might hold.

The remainder of this paper is organized as follows. In Section 2 we describe the model problem. In Section 3 we introduce the IETI-DP solver and we end this paper with numerical experiments in Section 4.

2 The problem setting

Let $\Omega \subset \mathbb{R}^2$ be open, simply connected and bounded with Lipschitz boundary $\partial\Omega$. $L_2(\Omega)$ and $H^1(\Omega)$ are the common Lebesgue and Sobolev spaces. As usual, $H_0^1(\Omega) \subset H^1(\Omega)$ denotes the subspace of functions that vanish on $\partial\Omega$.

We consider the following model problem: Find $u \in H_0^1(\Omega)$ such that

$$\int_{\Omega} \nabla u \cdot \nabla v \, dx = \int_{\Omega} f v \, dx \quad \text{for all } v \in H_0^1(\Omega) \quad (2)$$

with a given source function $f \in L_2(\Omega)$. We assume that Ω is a composition of K non-overlapping patches $\Omega^{(k)}$, where every patch $\Omega^{(k)}$ is parametrized by a geometry function

$$G_k : \widehat{\Omega} := (0, 1)^2 \rightarrow \Omega^{(k)} := G_k(\widehat{\Omega}) \subset \mathbb{R}^2,$$

that has a continuous extension to the closure of $\widehat{\Omega}$ and such that $\nabla G_k \in L_{\infty}(\widehat{\Omega})$ and $(\nabla G_k)^{-1} \in L_{\infty}(\widehat{\Omega})$.

We consider the case where the pre-images of the (Dirichlet) boundary consist of whole edges. The indices of neighboring patches $\Omega^{(\ell)}$ of $\Omega^{(k)}$, that share at least a part of their boundaries, is collected in the set

$$\mathcal{N}_{\Gamma}(k) := \{\ell \neq k : \text{meas}(\partial\Omega^{(k)} \cap \partial\Omega^{(\ell)}) > 0\},$$

where $\text{meas } T$ is the measure of T . For any $\ell \in \mathcal{N}_\Gamma(k)$, we write $\Gamma^{(k,\ell)} = \partial\Omega^{(k)} \cap \partial\Omega^{(\ell)}$. The endpoints of $\partial\Omega^{(k)} \cap \partial\Omega^{(\ell)}$ that are not located on the (Dirichlet) boundary of Ω are referred to as junctions. A junction could be a common vertex or a T-junction.

For the IgA discretization spaces, we first construct a B-spline space $\widehat{V}^{(k)}$ on the parameter domain $\widehat{\Omega}$ by tensorization of two univariate B-spline spaces. The function spaces on the physical domain are then defined by the pull-back principle: $V^{(k)} := \widehat{V}^{(k)} \circ G_k^{-1}$.

The product of the local spaces gives the global approximation space $V := V^{(1)} \times \dots \times V^{(K)}$. On this discretization space, we introduce the SIPG formulation, cf. [1, 15]. Since we are interested in a domain decomposition approach, we need patch-local formulations of SIPG.

3 The dG IETI-DP solver

For our patch-local formulations, we adapt the ideas of [3, 6, 5] and others. We choose local function spaces $V_e^{(k)}$ to be the product space of $V^{(k)}$ and the neighboring trace spaces $V^{(k,\ell)}$, which are the restrictions of $V^{(\ell)}$ to $\Gamma^{(k,\ell)}$. A function $v_e^{(k)} \in V_e^{(k)}$ is represented as a tuple $v_e^{(k)} = (v^{(k)}, (v^{(k,\ell)})_{\ell \in \mathcal{N}_\Gamma(k)})$, where $v^{(k)} \in V^{(k)}$ and $v^{(k,\ell)} \in V^{(k,\ell)}$. Note that the traces of the basis functions for $V^{(\ell)}$ restricted to $\Gamma^{(k,\ell)}$ form a basis of $V^{(k,\ell)}$. The basis for $V_e^{(k)}$ consists of the basis functions of $V^{(k)}$ and the basis functions for $V^{(k,\ell)}$. The basis functions on $V^{(k,\ell)}$ are usually visualized as living on artificial interfaces.

On each patch, we consider the local problem: Find $u_e^{(k)} \in V_e^{(k)}$ such that

$$\begin{aligned}
 a_e^{(k)}(u_e^{(k)}, v_e^{(k)}) &= \langle f_e^{(k)}, v_e^{(k)} \rangle \quad \text{for all } v_e^{(k)} \in V_e^{(k)}, \text{ where} \\
 a_e^{(k)}(u_e^{(k)}, v_e^{(k)}) &:= a^{(k)}(u_e^{(k)}, v_e^{(k)}) + m^{(k)}(u_e^{(k)}, v_e^{(k)}) + r^{(k)}(u_e^{(k)}, v_e^{(k)}), \\
 \langle f_e^{(k)}, v_e^{(k)} \rangle &:= \int_{\Omega^{(k)}} f v^{(k)} \, dx, \\
 a^{(k)}(u_e^{(k)}, v_e^{(k)}) &:= \int_{\Omega^{(k)}} \nabla u^{(k)} \cdot \nabla v^{(k)} \, dx, \\
 m^{(k)}(u_e^{(k)}, v_e^{(k)}) &:= \sum_{\ell \in \mathcal{N}_\Gamma(k)} \int_{\Gamma^{(k,\ell)}} \frac{\partial u^{(k)}}{\partial n_k} (v^{(k,\ell)} - v^{(k)}) \, ds, \\
 &\quad + \sum_{\ell \in \mathcal{N}_\Gamma(k)} \int_{\Gamma^{(k,\ell)}} \frac{\partial v^{(k)}}{\partial n_k} (u^{(k,\ell)} - u^{(k)}) \, ds, \\
 r^{(k)}(u_e^{(k)}, v_e^{(k)}) &:= \sum_{\ell \in \mathcal{N}_\Gamma(k)} \int_{\Gamma^{(k,\ell)}} \frac{\delta p^2}{h_{k\ell}} (u^{(k,\ell)} - u^{(k)}) (v^{(k,\ell)} - v^{(k)}) \, ds
 \end{aligned}$$

and n_k denotes the outward unit normal vector and δ is the dG penalty parameter, which has to be chosen large enough in order to guarantee that the bilinear form $a_e^{(k)}(\cdot, \cdot)$ is coercive. In [16], it was shown that δ can be chosen independently of p .

The discretization of $a_e^{(k)}(\cdot, \cdot)$ and $\langle f_e^{(k)}, \cdot \rangle$ gives a local system, which we write as

$$\begin{pmatrix} A_{\Pi}^{(k)} & A_{\Gamma}^{(k)} \\ A_{\Gamma}^{(k)} & A_{\Gamma\Gamma}^{(k)} \end{pmatrix} \begin{pmatrix} u_{\Pi}^{(k)} \\ u_{\Gamma}^{(k)} \end{pmatrix} = \begin{pmatrix} f_{\Pi}^{(k)} \\ f_{\Gamma}^{(k)} \end{pmatrix}, \tag{3}$$

where the index Π refers to the basis functions that are only supported in the interior of $\Omega^{(k)}$ and the index Γ refers to the remaining basis functions, i.e., those living on the patch boundary and on the artificial interfaces. We eliminate the interior degrees of freedom in (3) for every $k = 1, \dots, K$ to get the block diagonal Schur complement system

$$S\underline{w} = \underline{g},$$

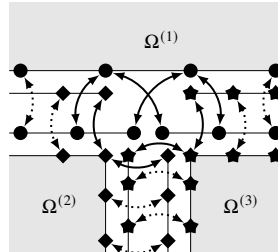
where the individual blocks of S are given by $S^{(k)} = A_{\Gamma\Gamma}^{(k)} - A_{\Gamma\Pi}^{(k)} (A_{\Pi\Pi}^{(k)})^{-1} A_{\Pi\Gamma}^{(k)}$.

The IETI-DP method requires carefully selected primal degrees of freedom to be solvable. We choose the degrees of freedom associated to the basis functions which are non-zero on a junction to be primal. For every standard corner, we only have one primal degree of freedom per patch, as in [15]. On a T-junction however, the number of non-zero basis functions grows linearly with p . Since we take all of them, we refer to “fat vertices” in this context.

$C = \text{diag}(C^{(1)}, \dots, C^{(K)})$ is the constraint matrix, i.e., it is defined such that $C\underline{w} = 0$ if and only if the associated function w vanishes at the primal degrees of freedom. The matrix Ψ represents the energy minimizing basis functions for the space of primal degrees of freedom.

Furthermore, we introduce the jump matrix B , which models the jumps of the functions between the patch boundaries and the associated artificial interfaces. Each row corresponds to one degree of freedom (coefficient for a basis function) on the the patch boundary and one artificial interface; as usual, each row has only two non-zero coefficients that are -1 and 1 . Primal degrees of freedom are excluded. For a visualization, see Fig. 1, where the primal degrees of freedom are marked with solid lines and the dotted arrows show the action of the jump matrix B . The basis functions on the artificial interfaces are labeled with the same symbols from the original spaces.

Fig. 1: Action of matrix B (dotted lines) and primal degrees of freedom (solid lines)



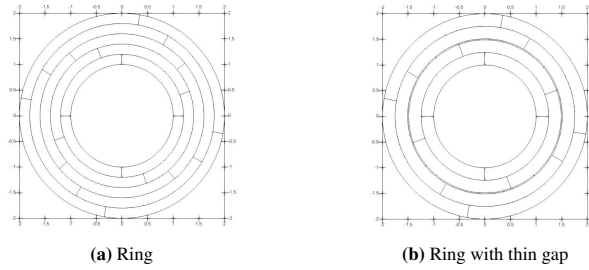


Fig. 2: Computational domains and the decomposition into patches

patches in Fig. 2a, the condition number grows only by a factor between 5 and 6. Also the iteration counts grow only mildly.

r	Fig. 2a								Fig. 2b							
	p = 2		p = 3		p = 6		p = 7		p = 2		p = 3		p = 6		p = 7	
	it	κ	it	κ	it	κ	it	κ	it	κ	it	κ	it	κ	it	κ
4	9	3.7	9	3.5	8	2.4	8	2.1	12	18.0	13	18.0	13	12.9	11	11.6
5	10	4.6	10	4.5	9	3.8	9	3.5	20	24.1	19	23.4	19	19.3	18	18.0
6	10	5.8	10	5.5	10	4.9	10	4.8	22	31.7	22	29.9	21	25.9	20	24.8
7	11	6.3	11	6.2	10	5.6	10	5.5	24	37.2	24	36.3	22	31.3	22	30.1
8	11	6.7	11	6.7	11	6.3	10	5.6	24	43.2	24	42.3	24	36.5	24	31.6

Table 1: Iterations it and condition numbers κ ; degree p ; refinement level r

The Table 2 presents the parallel solving times for n processors. We only consider the domain in Fig. 2a again with $s = p - 1$. We see that the speedup rate with respect to n is a bit smaller than the expected rate of 2. This is probably caused by the rather small number of patches in the computational domain. In Table 3 we report on the iteration counts and the condition numbers for the decomposition in Fig. 2a when we change the smoothness s of the B-splines within the patches. The numbers in the table show the behavior for $r = 5$. We see that for a fixed smoothness s the condition number grows slightly with respect to the spline degree p . For a fixed degree p , we observe a decline in the condition number when we increase the smoothness s .

Acknowledgements The first author was supported by the Austrian Science Fund (FWF): S117 and W1214-04. The second author has also received support from the Austrian Science Fund (FWF): P31048.

r	$p = 3$					$p = 7$				
	$n = 1$	$n = 2$	$n = 4$	$n = 8$	$n = 16$	$n = 1$	$n = 2$	$n = 4$	$n = 8$	$n = 16$
6	3.8	2.8	2.4	1.2	0.8	10.0	6.5	5.0	2.5	1.75
7	24.0	16.1	13.6	6.4	4.1	47.0	31.7	26.1	12.2	9.3
8	107.0	81.4	66.8	29.5	19.5	220.0	158.7	129.4	56.7	45.4

Table 2: Solving times (sec.); degree p ; refinement level r ; n processors; Fig. 2a

s	$p = 2$		$p = 3$		$p = 4$		$p = 5$		$p = 6$		$p = 7$	
	it	κ	it	κ	it	κ	it	κ	it	κ	it	κ
0	10	5.0	10	5.3	10	5.4	10	5.5	10	5.6	10	5.6
1	10	4.6	10	5.2	10	5.3	10	5.4	10	5.5	10	5.5
2			10	4.5	10	5.0	10	5.3	10	5.4	10	5.5
3					9	4.2	10	4.9	10	5.1	10	5.3
4							9	4.0	10	4.7	10	5.0
5									9	3.8	9	4.5
6											9	3.5

Table 3: Iterations it and condition number κ ; refinement level $r = 5$; degree p ; smoothness s ; Fig. 2a

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